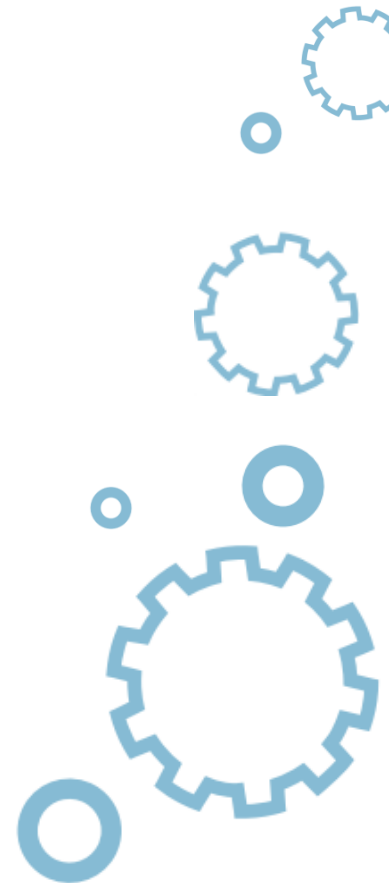




Maciej Sobieraj

Lecture 12



Outline

1. Linked Lists

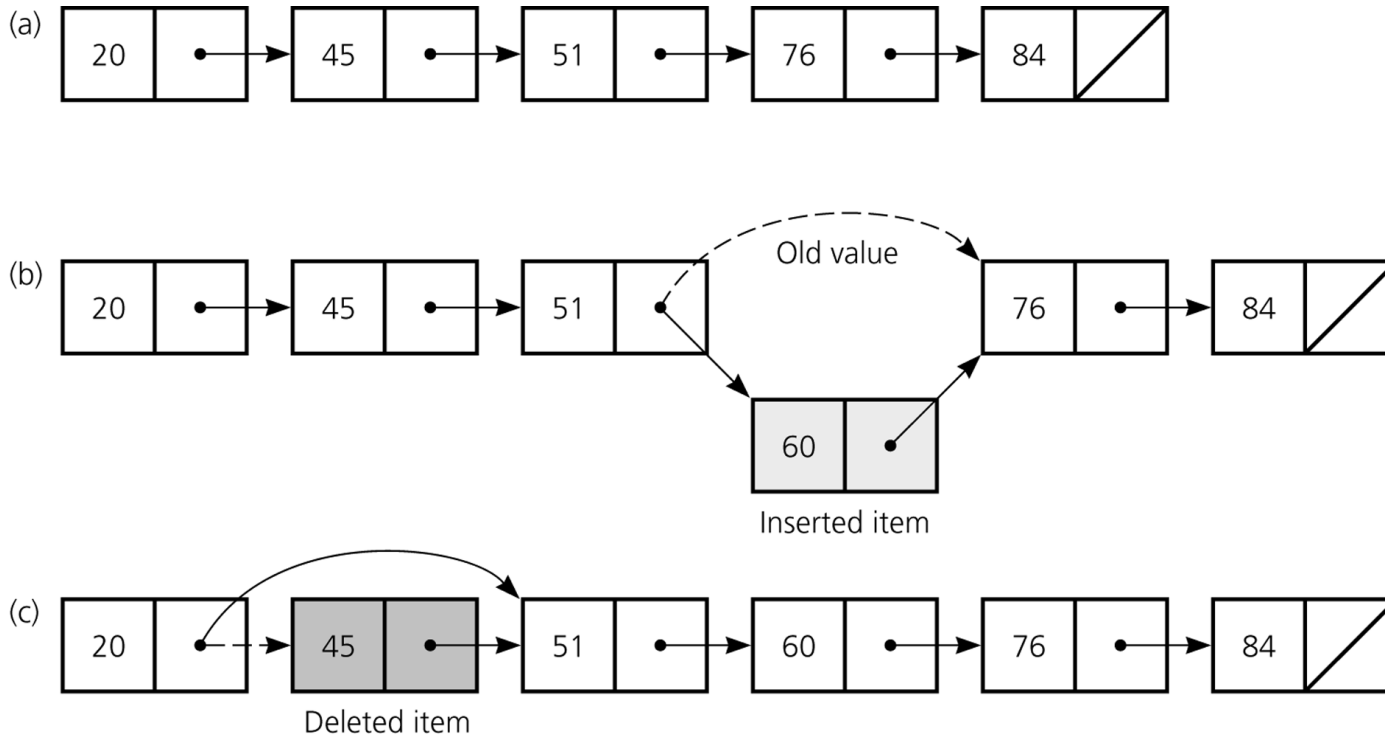


Preliminaries

- Options for implementing an ADT List
 - **Array has a fixed size**
 - Data must be shifted during insertions and deletions
 - **Linked list is able to grow in size as needed**
 - Does not require the shifting of items during insertions and deletions



Preliminaries



a) A linked list of integers; b) insertion; c) deletion



Pointers

- A pointer contains the location, or address in memory, of a memory cell
 - Initially undefined, but not *NULL*
 - A statically allocated pointer declaration

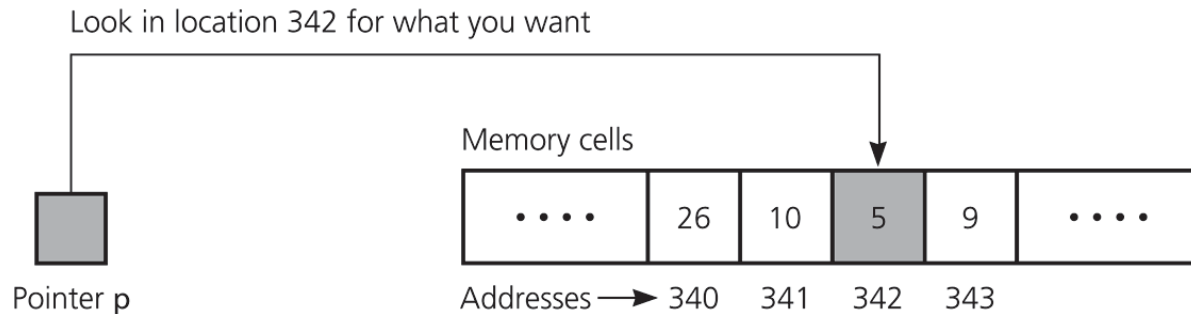
```
int *p;
```
 - A dynamically allocated pointer variable

```
p = new int;
```



Pointers

- The expression, $*p$, denotes the memory cell to which p points
- The $\&$ address-of operator places the address of a variable into a pointer variable



A pointer to an integer



Pointers

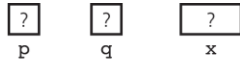
- The `delete` operator returns dynamically allocated memory to the system for reuse, and leaves the variable undefined
 - `delete p;`
 - A pointer to a deallocated memory cell is possible and dangerous
- Assign the pointer `q` the value in `p`

```
q = p;
```

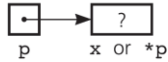


Pointers

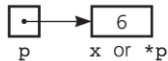
(a) `int *p, *q;`
`int x;`



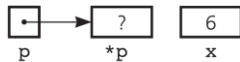
(b) `p = &x;`



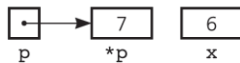
(c) `*p = 6;`



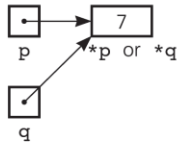
(d) `p = new int;`



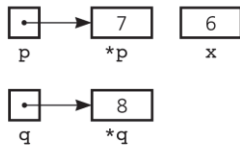
(e) `*p = 7;`



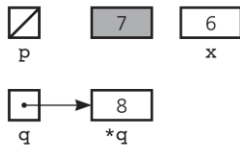
(f) `q = p;`



(g) `q = new int;`
`*q = 8;`



(h) `p = NULL;`



(i) `delete q;`
`q = NULL;`

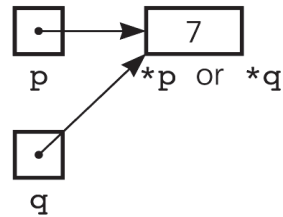


- (a) declaring pointer variables;
- (b) pointing to statically allocating memory;
- (c) assigning a value;
- (d) allocating memory dynamically;
- (e) assigning a value

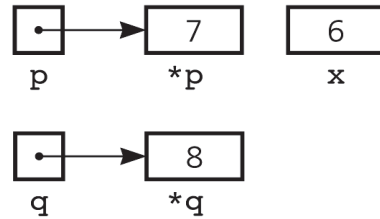


Pointers

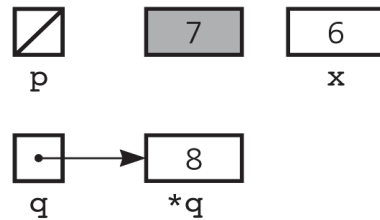
(f) `q = p;`



(g) `q = new int;`
`*q = 8;`



(h) `p = NULL;`



(i) `delete q;`
`q = NULL;`



(f) copying a pointer;

(g) allocating memory dynamically
and assigning a value;

(h) assigning NULL to a pointer variable;

(i) deallocating memory



Dynamic Allocation of Arrays

- Use the `new` operator to allocate an array dynamically
- An array name is a pointer to the array's first element
- The size of a dynamically allocated array can be increased

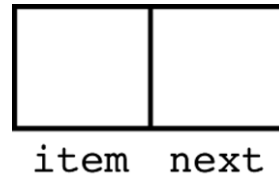
```
double* oldArray = anArray;  
anArray = new double [2*arraySize];
```



Pointer-Based Linked Lists

- A node in a linked list is usually a `struct`

```
struct Node
{ int item
  Node *next;
}; //end struct
```



A node

- A node is dynamically allocated

```
Node *p;
p = new Node;
```

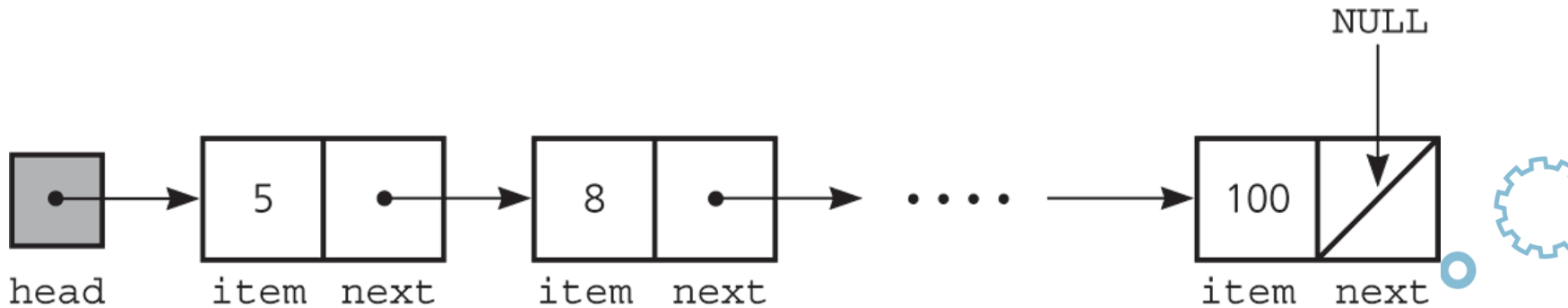


Pointer-Based Linked Lists

- The head pointer points to the first node in a linked list
- If head is *NULL*, the linked list is empty
- Executing the statement `head=new Node` before `head=NULL` will result in a lost cell



Pointer-Based Linked Lists



A head pointer to a list



A lost cell



Displaying the Contents of a Linked List

- Reference a node member with the `->` operator

```
p->item;
```

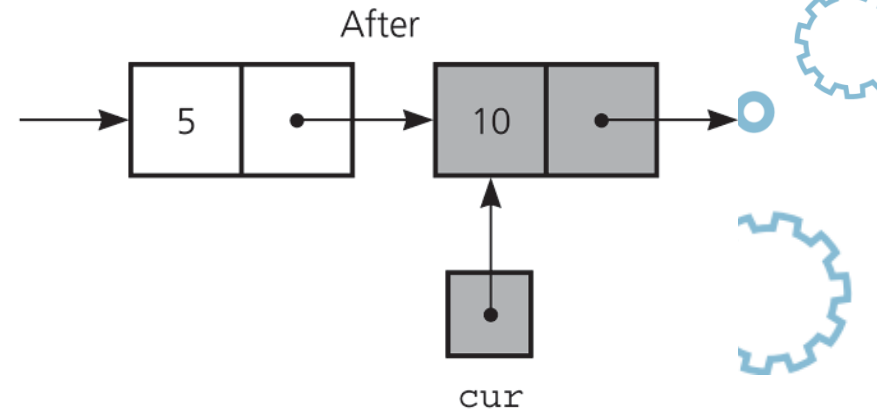
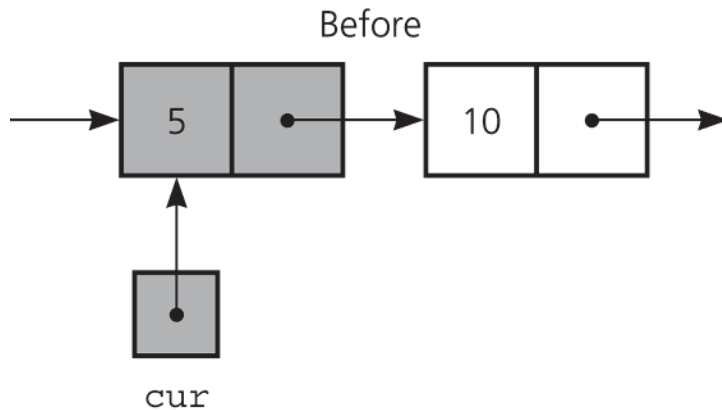
- A traverse operation visits each node in the linked list

- A pointer variable `cur` keeps track of the current node

```
for (Node *cur = head;  
     cur != NULL; cur = cur->next)  
    cout << cur->item << endl;
```



Displaying the Contents of a Linked List



The effect of the assignment $cur = cur \rightarrow next$



Deleting a Specified Node from a Linked List

- Deleting an interior node

```
prev->next=cur->next;
```

- Deleting the first node

```
head=head->next;
```

- Return deleted node to system

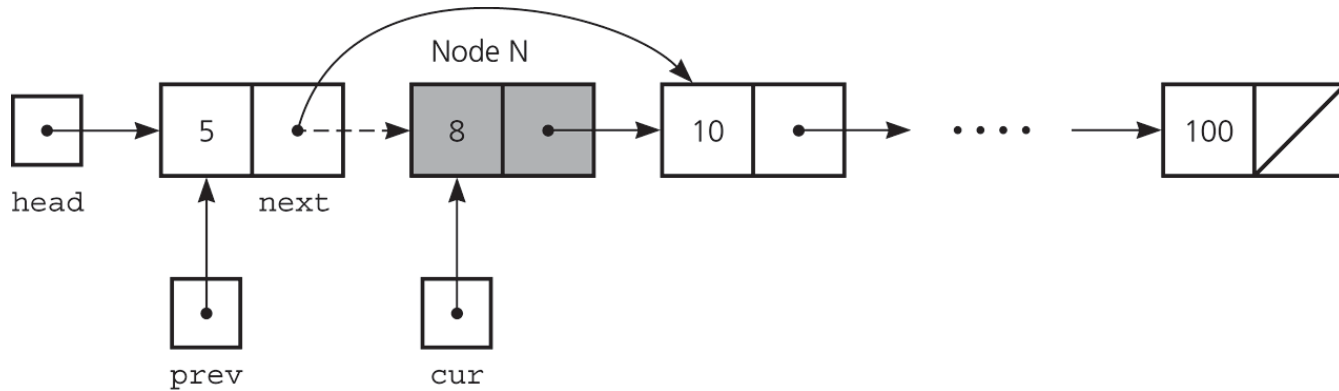
```
cur->next = NULL;
```

```
delete cur;
```

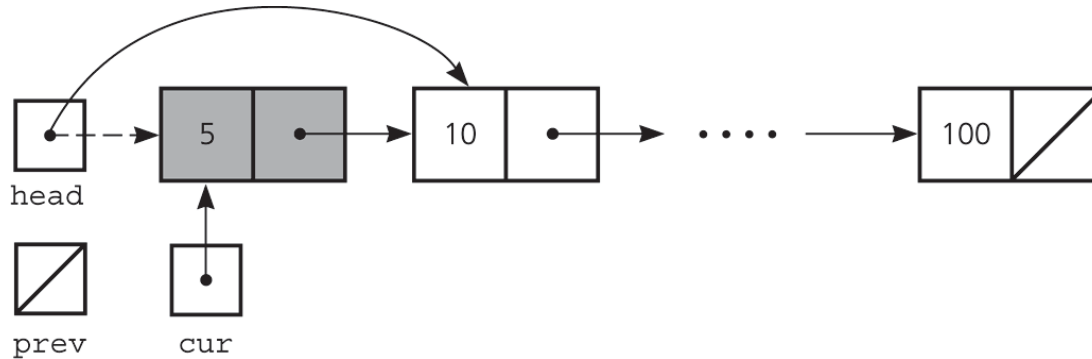
```
cur=NULL;
```



Deleting a Specified Node from a Linked List



Deleting a node from a linked list



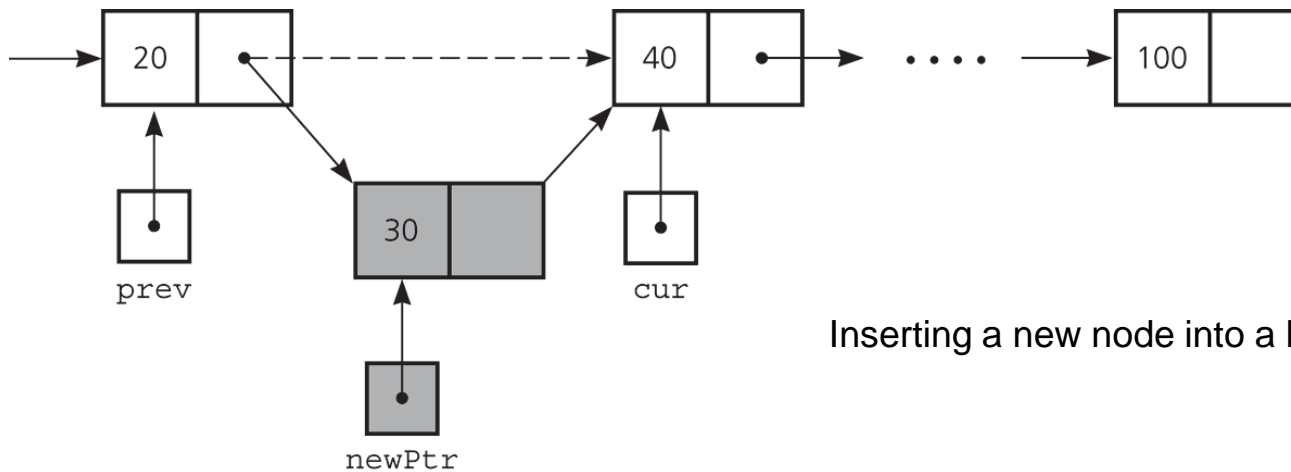
Deleting the first node



Inserting a Node into a Specified Position of a Linked List

- To insert a node between two nodes

```
newPtr->next = cur;  
prev->next = newPtr;
```



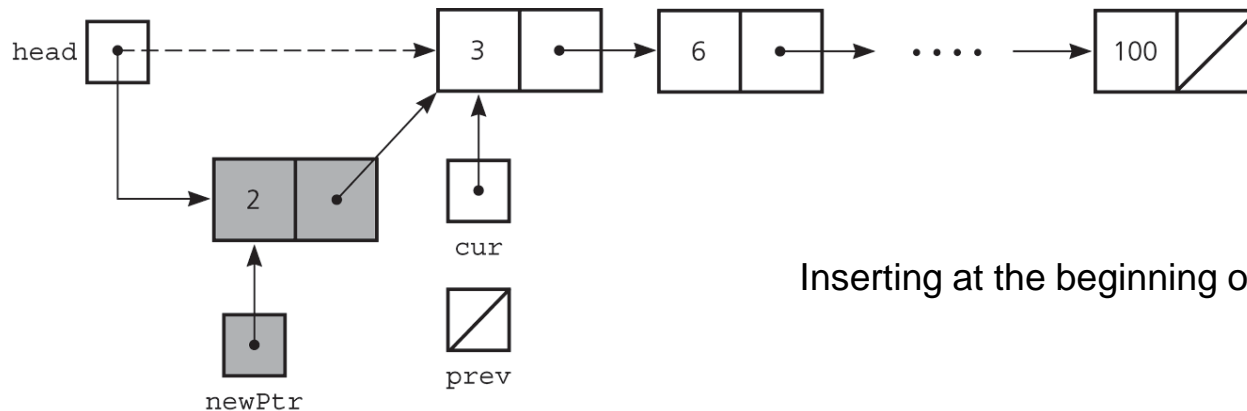
Inserting a new node into a linked list



Inserting a Node into a Specified Position of a Linked List

- To insert a node at the beginning of a linked list

```
newPtr->next = head;  
head = newPtr;
```



Inserting at the beginning of a linked list

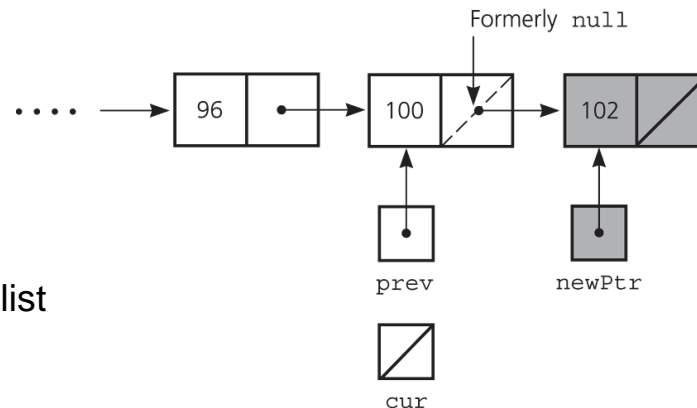


Inserting a Node into a Specified Position of a Linked List

- Inserting at the end of a linked list is not a special case if `cur` is `NULL`

```
newPtr->next = cur;
```

```
prev->next = newPtr;
```



Inserting at the end of a linked list



Inserting a Node into a Specified Position of a Linked List

- Determining the point of insertion or deletion for a sorted linked list of objects

```
for (prev = NULL, cur= head;  
      (cur != null) &&  
      (newValue > cur->item);  
      prev = cur, cur = cur->next;
```



A Pointer-Based Implementation of the ADT List

- Public methods

- `isEmpty`
- `getLength`
- `insert`
- `remove`
- `retrieve`

- Private method

- `find`

- Private Data

Members

- `head`
- `Size`
- Local variables to member functions
 - `cur`
 - `prev`



Constructors and Destructors

- Default constructor initializes size and head
- Copy constructor allows a deep copy
 - Copies the array of list items and the number of items
- A destructor is required for dynamically allocated memory



Comparing Array-Based and Pointer-Based Implementations

- Size
 - Increasing the size of a resizable array can waste storage and time
- Storage requirements
 - Array-based implementations require less memory than a pointer-based ones



Comparing Array-Based and Pointer-Based Implementations

- Access time
 - Array-based: constant access time
 - Pointer-based: the time to access the i^{th} node depends on i
- Insertion and deletions
 - Array-based: require shifting of data
 - Pointer-based: require a list traversal



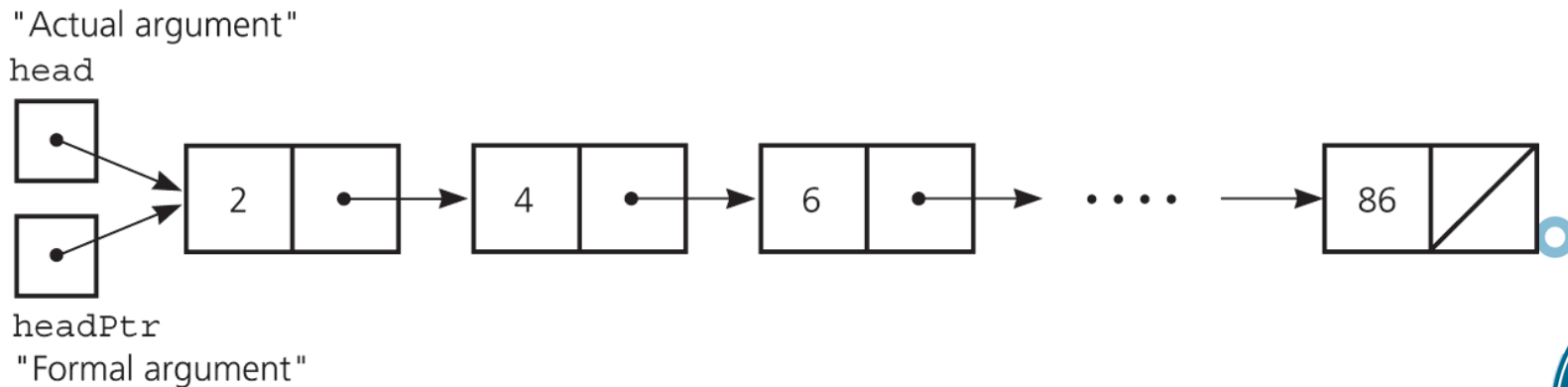
Saving and Restoring a Linked List by Using a File

- Use an external file to preserve the list between runs
- Do not write pointers to a file, only data
- Recreate the list from the file by placing each item at the end of the list
 - Use a tail pointer to facilitate adding nodes to the end of the list
 - Treat the first insertion as a special case by setting the tail to head



Passing a Linked List to a Function

- A function with access to a linked list's `head` pointer has access to the entire list
- Pass the head pointer to a function as a reference argument



A head pointer as a value argument



Processing Linked Lists Recursively

- Recursive strategy to display a list
 - Write the first node of the list
 - Write the list minus its first node
- Recursive strategies to display a list backward
 - `writeListBackward` strategy
 - Write the last node of the list
 - Write the list minus its last node backward



Processing Linked Lists Recursively

- `writeListBackward2` strategy
 - Write the list minus its first node backward
 - Write the first node of the list
- Recursive view of a sorted linked list
 - The linked list to which `head` points is a sorted list if
 - `head` is *NULL* or
 - `head->next` is *NULL* or
 - `head->item < head->next->item`, and `head->next` points to a sorted linked list



Objects as Linked List Data

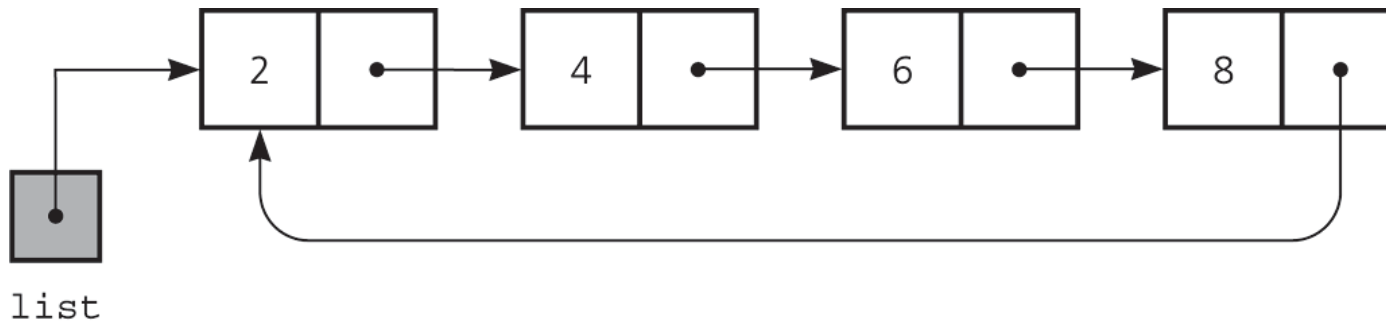
- Data in a linked list node can be an instance of a class

```
typedef ClassName ItemType;  
struct Node  
{ ItemType item;  
  Node *next;  
}; //end struct  
Node *head;
```



Circular Linked Lists

- Last node references the first node
- Every node has a successor
- No node in a circular linked list contains *NULL*



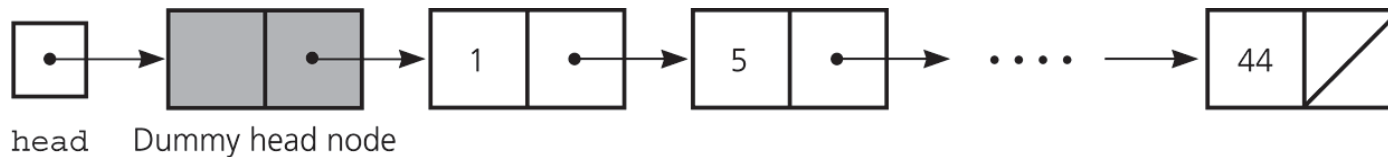
A circular linked list



Dummy Head Nodes

- Dummy head node

- Always present, even when the linked list is empty
- Insertion and deletion algorithms initialize `prev` to reference the dummy head node, rather than `NULL`



A dummy head node



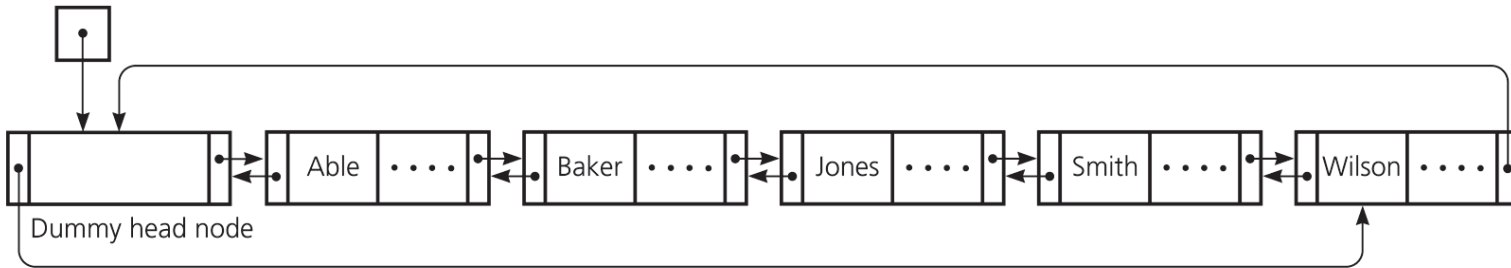
Doubly Linked Lists

- Each node points to both its predecessor and its successor
- Circular doubly linked list
 - `precede` pointer of the dummy head node points to the last node
 - `next` reference of the last node points to the dummy head node
 - No special cases for insertions and deletions

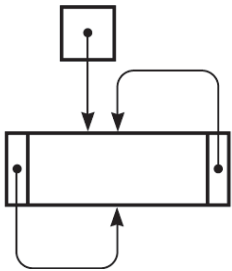


Doubly Linked Lists

(a) listHead



(b) listHead



(a) A circular doubly linked list with a dummy head node

(b) An empty list with a dummy head node



Doubly Linked Lists

- To delete the node to which `cur` points

```
(cur->precede)->next = cur->next;  
(cur->next)->precede = cur->precede;
```
- To insert a new node pointed to by `newPtr` before the node pointed to by `cur`

```
newPtr->next = cur;  
newPtr->precede = cur->precede;  
cur->precede = newPtr;  
newPtr->precede->next = newPtr;
```



Application: Maintaining an Inventory

- Operations on the inventory
 - List the inventory in alphabetical order by title (L command)
 - Find the inventory item associated with title (I, M, D, O, and S commands)
 - Replace the inventory item associated with a title (M, D, R, and S commands)
 - Insert new inventory items (A and D commands)



The C++ Standard Template Library

- The STL contains class templates for some common ADTs, including the *list* class
- The STL provides support for predefined ADTs through three basic items
 - Containers are objects that hold other objects
 - Algorithms act on containers
 - Iterators provide a way to cycle through the contents of a container



Summary

- The C++ `new` and `delete` operators enable memory to be dynamically allocated and recycled
- Each pointer in a linked list is a pointer to the next node in the list
- Array-based lists use an implicit ordering scheme; pointer-based lists use an explicit ordering scheme



Summary

- Algorithms for insertions and deletions in a linked list involve traversing the list and performing pointer changes
 - Inserting a node at the beginning of a list and deleting the first node of a list are special cases
- A class that allocates memory dynamically needs an explicit copy constructor and destructor



Summary

- Recursion can be used to perform operations on a linked list
- In a circular linked list, the last node points to the first node
- Dummy head nodes eliminate the special cases for insertion into and deletion from the beginning of a linked list

